

# EN164: Design of Computing Systems

## Lecture 11: Processor / ISA 4

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Spring 2011



[ material from Patterson & Hennessy, 4<sup>th</sup> ed and Harris 1<sup>st</sup> ed ]

# Procedure calls

## Procedure calling conventions:

- Caller:
  - passes **arguments** to callee.
  - jumps to the callee
- Callee:
  - **performs the procedure**
  - **returns the result** to caller
  - **returns to the point of call**
  - **must not overwrite** registers or memory needed by the caller

## MIPS conventions:

- Call procedure: jump and link ( $j\ a\ l$ )
- Return from procedure: jump register ( $j\ r$ )
- Argument values:  $\$a0 - \$a3$
- Return value:  $\$v0 - \$v1$

# Procedure calls

## High-level code

```
int main() {
    simple();
    a = b + c;
}

void simple() {
    return;
}
```

## MIPS assembly code

```
0x00400200 main: jal  simple
0x00400204          add  $s0, $s1, $s2
...

0x00401020 simple: jr  $ra
```

`void` means that `simple` doesn't return a value.

`jal`: jumps to `simple` and saves PC+4 in the return address register (`$ra`).  
In this case, `$ra = 0x00400204` after `jal` executes.

`jr $ra`: jumps to address in `$ra`, in this case `0x00400204`.

# Input arguments and return values

## High-level code

```
int main()
{
    int y;
    ...
    y = diffofsums(2, 3, 4, 5); // 4 arguments
    ...
}

int diffofsums(int f, int g, int h, int i)
{
    int result;
    result = (f + g) - (h + i);
    return result;           // return value
}
```

# Input arguments and return values

## MIPS assembly code

main:

...

```
addi $a0, $0, 2    # argument 0 = 2
addi $a1, $0, 3    # argument 1 = 3
addi $a2, $0, 4    # argument 2 = 4
addi $a3, $0, 5    # argument 3 = 5
jal  diffofsums    # call procedure
add  $s0, $v0, $0  # y = returned value
```

...

# \$s0 = result

diffofsums:

```
add $t0, $a0, $a1  # $t0 = f + g
add $t1, $a2, $a3  # $t1 = h + i
sub $s0, $t0, $t1  # result = (f + g) - (h + i)
add $v0, $s0, $0   # put return value in $v0
jr  $ra           # return to caller
```

## MIPS conventions:

- Argument values: \$a0 - \$a3
- Return value: \$v0

# Potential problems in procedural calling

## MIPS assembly code

```
# $s0 = result
diffofsums:
    add $t0, $a0, $a1 # $t0 = f + g
    add $t1, $a2, $a3 # $t1 = h + i
    sub $s0, $t0, $t1 # result = (f + g) - (h + i)
    add $v0, $s0, $0 # put return value in $v0
    jr $ra # return to caller
```

- `diffofsums` overwrote 3 registers: `$t0`, `$t1`, and `$s0`
- `diffofsums` can use the *stack* to temporarily store registers

# The stack

- Memory used to temporarily save variables
- Like a stack of dishes, last-in-first-out (LIFO) queue
- *Expands*: uses more memory by growing down (from higher to lower memory addresses) when more space is needed
- *Contracts*: uses less memory when the space is no longer needed
- Stack pointer:  $\$sp$ , points to top of the stack



Address	Data		Address	Data	
7FFFFFFC	12345678	← \$sp	7FFFFFFC	12345678	
7FFFFFF8			7FFFFFF8	AABBCCDD	
7FFFFFF4			7FFFFFF4	11223344	← \$sp
7FFFFFF0			7FFFFFF0		
⋮	⋮		⋮	⋮	
⋮	⋮		⋮	⋮	

# How procedures use the stack

- Called procedures must have no other unintended side effects.
- But `diffofsums` overwrites 3 registers: `$t0`, `$t1`, `$s0`

```
# MIPS assembly
# $s0 = result
diffofsums:
    add $t0, $a0, $a1    # $t0 = f + g
    add $t1, $a2, $a3    # $t1 = h + i
    sub $s0, $t0, $t1    # result = (f + g) - (h + i)
    add $v0, $s0, $0     # put return value in $v0
    jr  $ra              # return to caller
```



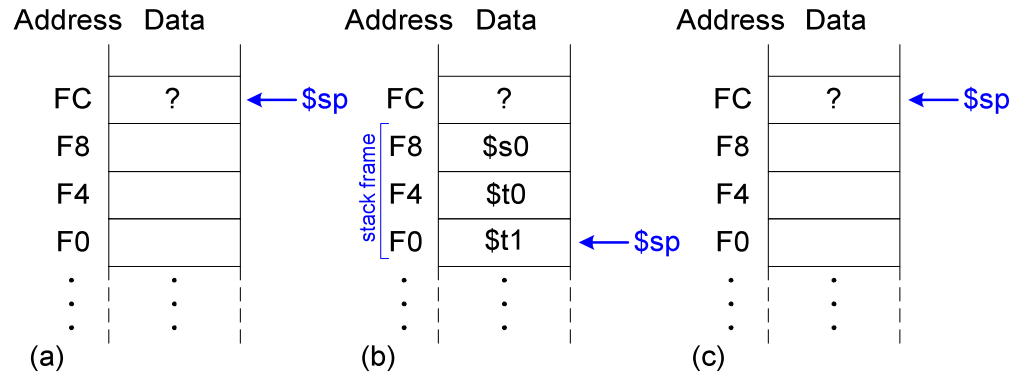
# Storing register values on the stack

```

# $s0 = result
diffofsums:
    addi $sp, $sp, -12    # make space on stack
                          # to store 3 registers

    sw   $s0, 8($sp)     # save $s0 on stack
    sw   $t0, 4($sp)     # save $t0 on stack
    sw   $t1, 0($sp)     # save $t1 on stack
    add  $t0, $a0, $a1    # $t0 = f + g
    add  $t1, $a2, $a3    # $t1 = h + i
    sub  $s0, $t0, $t1    # result = (f + g) - (h + i)
    add  $v0, $s0, $0     # put return value in $v0
    lw   $t1, 0($sp)     # restore $t1 from stack
    lw   $t0, 4($sp)     # restore $t0 from stack
    lw   $s0, 8($sp)     # restore $s0 from stack
    addi $sp, $sp, 12    # deallocate stack space
    jr   $ra             # return to caller

```



# Protocol for preserving registers

Preserved <i>Callee-Saved</i>	Nonpreserved <i>Caller-Saved</i>
<code>\$s0 - \$s7</code>	<code>\$t0 - \$t9</code>
<code>\$ra</code>	<code>\$a0 - \$a3</code>
<code>\$sp</code>	<code>\$v0 - \$v1</code>
stack above <code>\$sp</code>	stack below <code>\$sp</code>

# Multiple procedure calls

```
proc1:  
    addi $sp, $sp, -4    # make space on stack  
    sw   $ra, 0($sp)    # save $ra on stack  
    jal  proc2  
    ...  
    lw   $ra, 0($sp)    # restore $s0 from stack  
    addi $sp, $sp, 4    # deallocate stack space  
    jr   $ra            # return to caller
```

# Storing saved registers on the stack

```
# $s0 = result
diffofsums:
    addi $sp, $sp, -4    # make space on stack to
                        # store one register
    sw  $s0, 0($sp)    # save $s0 on stack
                        # no need to save $t0 or $t1
    add $t0, $a0, $a1    # $t0 = f + g
    add $t1, $a2, $a3    # $t1 = h + i
    sub $s0, $t0, $t1    # result = (f + g) - (h + i)
    add $v0, $s0, $0     # put return value in $v0
    lw  $s0, 0($sp)    # restore $s0 from stack
    addi $sp, $sp, 4    # deallocate stack space
    jr  $ra              # return to caller
```

# Recursive procedure calls

## High-level code

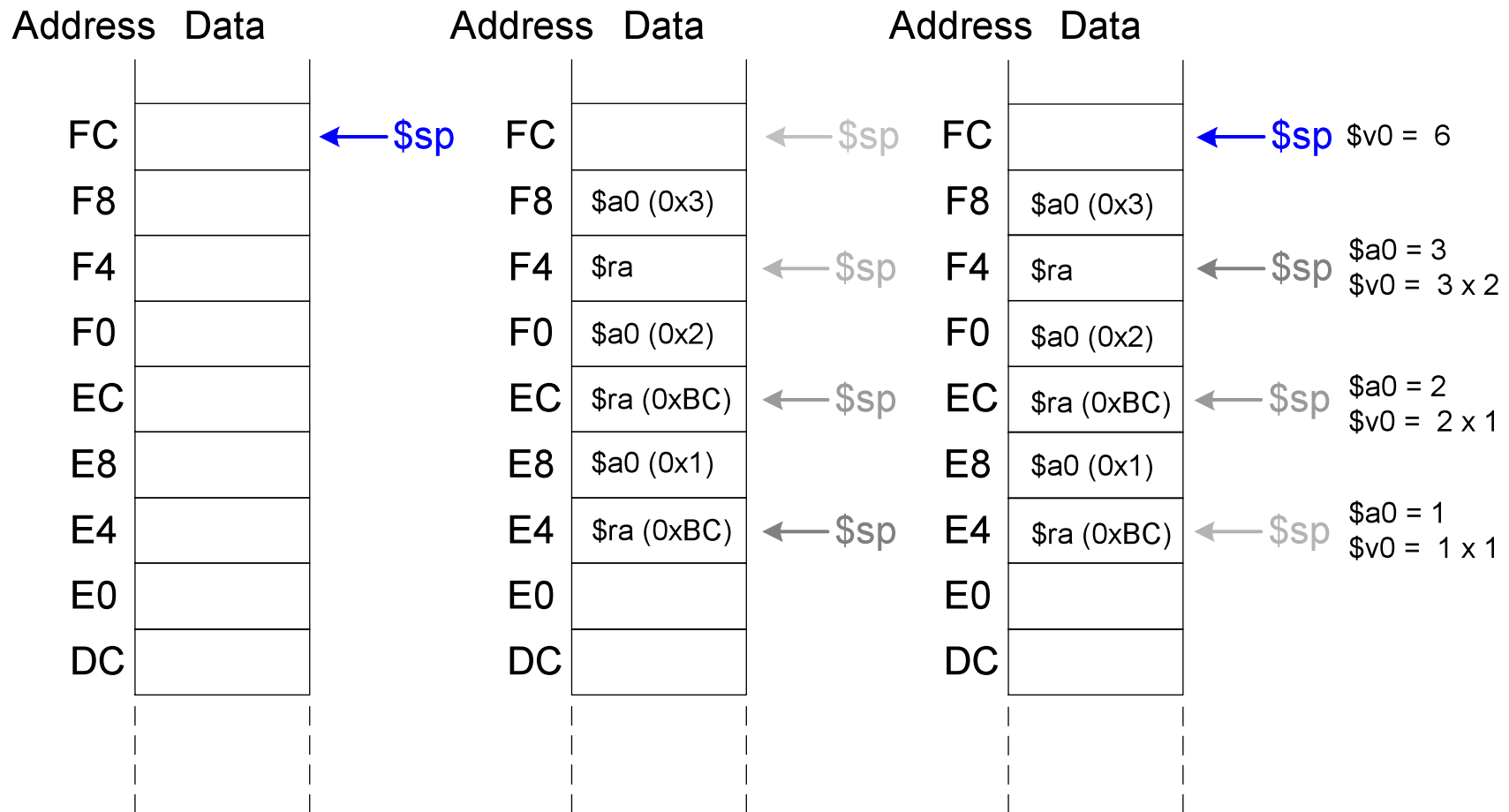
```
int factorial(int n) {  
    if (n <= 1)  
        return 1;  
    else  
        return (n * factorial(n-1));  
}
```

# Recursive procedure calls

## MIPS assembly code

```
0x90 factorial: addi $sp, $sp, -8 # make room
0x94             sw  $a0, 4($sp) # store $a0
0x98             sw  $ra, 0($sp) # store $ra
0x9C             addi $t0, $0, 2
0xA0             slt  $t0, $a0, $t0 # a <= 1 ?
0xA4             beq  $t0, $0, else # no: go to else
0xA8             addi $v0, $0, 1 # yes: return 1
0xAC             addi $sp, $sp, 8 # restore $sp
0xB0             jr   $ra # return
0xB4             else: addi $a0, $a0, -1 # n = n - 1
0xB8             jal  factorial # recursive call
0xBC             lw   $ra, 0($sp) # restore $ra
0xC0             lw   $a0, 4($sp) # restore $a0
0xC4             addi $sp, $sp, 8 # restore $sp
0xC8             mul  $v0, $a0, $v0 # n * factorial(n-1)
0xCC             jr   $ra # return
```

# Stack during recursive calls



# Procedural call summary

- Caller
  - Put arguments in \$a0-\$a3
  - Save any registers that are needed (\$ra, maybe \$t0-t9)
  - jal callee
  - Restore registers
  - Look for result in \$v0
- Callee
  - Save registers that might be disturbed (\$s0-\$s7)
  - Perform procedure
  - Put result in \$v0
  - Restore registers
  - jr \$ra